

# $\Delta$ -coloring in the graph streaming model

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  - ▶ Median-finding ( $O(p)$  pass  $\tilde{O}(n^{1/p})$  space algorithms are known).
  - ▶ Counting distinct elements ( $(1 + \epsilon)$ -approx in  $O(1/\epsilon^2 \cdot \log n)$  space with constant probability of success known).

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- ▶ Semi-streaming model: Allow  $O(n \cdot \text{polylog}(n))$  space.
- ▶ A good survey is [2]

## $\Delta$ -coloring

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- ▶ For  $G$  with maximum degree  $\Delta$ ,  $\Delta + 1$  coloring is easy (assign arb. color to some vertex, then color rest “greedily”).
- ▶ If  $G$  is not a clique or an odd-cycle, then it can be  $\Delta$ -colored (Brook’s Theorem).

# Our project

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- ▶ It is known that randomized  $(\Delta + 1)$ -coloring is possible in  $\tilde{O}(n)$  space [1]
- ▶ With randomization, can we  $\Delta$ -color graphs in  $\tilde{O}(n)$  space in  $O(1)$  passes?



# References



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