

# Creating and Classifying an Alternate Hierarchy

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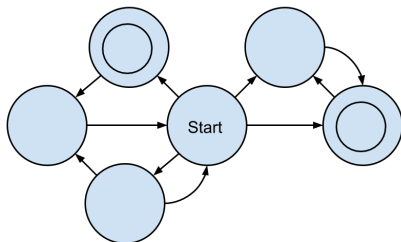
Harvey Mudd College

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# Deterministic Polynomial Time

## Definition

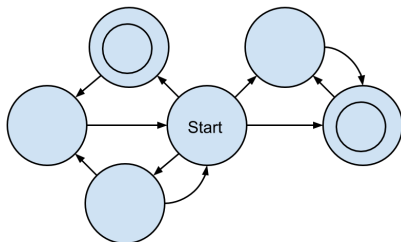
A language  $L$  is in  $\mathbf{P}$  if for any  $x \in L$ , there is a polynomial time machine  $M$  such that  $M(x) = 1$ .



# Nondeterministic Polynomial Time

## Definition

A language  $L$  is in **NP** if for any  $x \in L$ , there is a polynomial time machine  $M$  such that  $\exists a M(x, a) = 1$ .



# Nondeterministic Polynomial Time (Complement)

## Definition

A language  $L$  is in **coNP** if for any  $x \in L$ , there is a polynomial time machine  $M$  such that  $\forall a M(x, a) = 1$ .

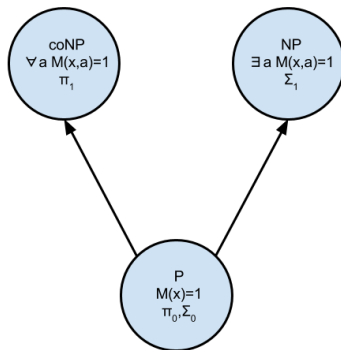
Note, the complement of an NP problem is:

$$\neg \exists a M_1(x, a) = 1,$$

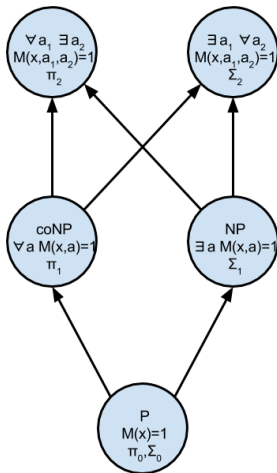
$$\forall a M_1(x, a) = 0,$$

$$\forall a M_2(x, a) = 1.$$

# Polynomial Hierarchy

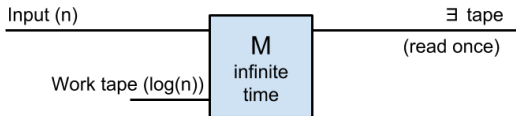
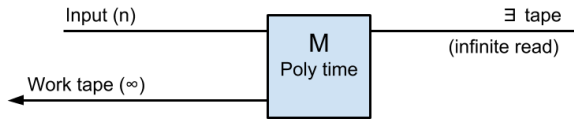


# Polynomial Hierarchy

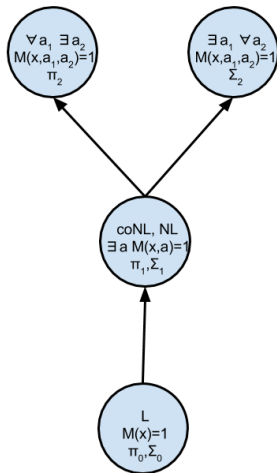


# Defining a New Logspace Hierarchy

Polynomial machine vs our machine

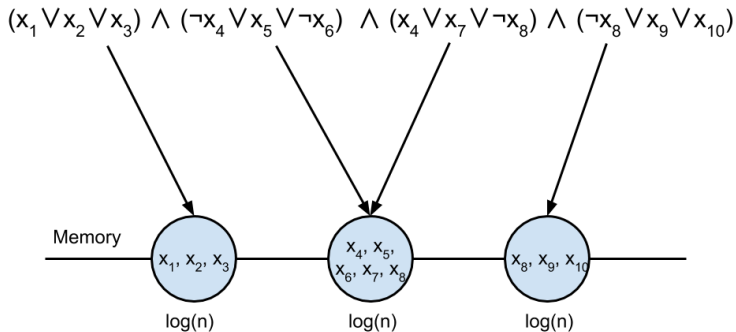


# Defining a New Logspace Hierarchy





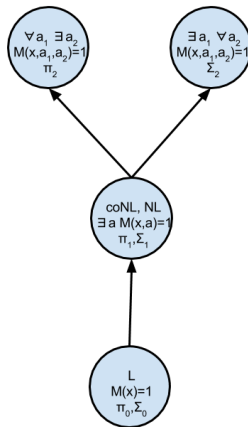
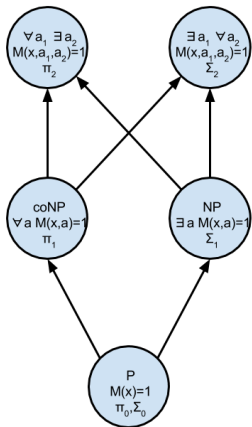
# The Pathwidth Problem



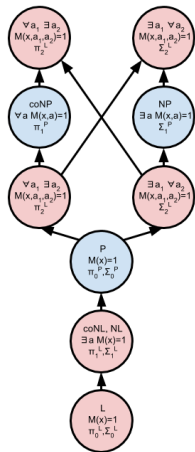
# $\hat{\Sigma}_1^L$ Complete Problem

- ▶ Proof of completeness
  - ▶ Containment: clearly,  $M$  can solve bounded pathwidth  $SAT$
  - ▶ Hardness: apply Cook-Levin Theorem to represent  $M$  as boolean expression, each step uses only logspace, so the expression is a bounded pathwidth  $SAT$  problem
  
- ▶ This works for other levels of the hierarchy
  - ▶ ie,  $\forall a_1 \exists a_2 SAT$  with logarithmically bounded pathwidth is complete for  $\hat{\Pi}_2^L$

# Comparisons with the Polynomial Hierarchy



# Combined Hierarchy



# Conclusion



- ▶ We've created a new machine that creates a hierarchy that probably doesn't collapse
  - ▶ Else the polynomial hierarchy would collapse!
- ▶ We know a characterization for the complete problems for the entire hierarchy
- ▶ We know how it fits with the polynomial hierarchy
- ▶ But it isn't helpful to the Group Isomorphism problem (see: Czechs students' presentation)
- ▶ We hope to find more natural complete problems in the hierarchy

# Acknowledgments

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# Bibliography

-  Sanjeev Arora and Boaz Barak: *Computational Complexity: A Modern Approach*. Cambridge University Press, 2007.
-  Michael Garey and David Johnson: *Computers and Intractability: A Guide to the Theory of NP-Completeness*. WH Freemanx, 1979.