# Designing non-manipulable tournament rules 

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## Introduction: an example

- Given $\binom{n}{2}$ outcomes between $n$ teams (= tournament), determine the winner of the tournament (= tournament rule)
- Example: random single binary elimination bracket (RS2EB)



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$\{A, D\}$-adjecent tournament wrt $T$

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- We say that the rule RS2EB is $2-S N M-\frac{1}{3}$


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- (CC) If a team $A$ beats every other team, then the chance that $A$ is the winner of $T$ is 1


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There is a rule that is both (CC) and ( $k-S N M-\frac{2}{3}$ ) but not monotone.

Theorem (Dinev and Weinberg, 2022; not published yet)
There exists a rule that is monotone, (CC), (2-SNM- $\frac{1}{3}$ ), and (3-SNM- $\frac{31}{60}$ ).

## Random single ternary elimination bracket (RS3EB)



Theorem
The rule RS3EB is monotone, (CC), and (3-SNM- $\alpha$ ) for $\frac{31}{60}<\frac{227}{420} \leq \alpha \leq \frac{23}{27}$.

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- We have only two main restrictions imposed on these rules
(CC) condition
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$$
r_{a}(T)=1
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( $k$-SNM- $\alpha$ ) condition

- Impacts only specific pairs of tournaments



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- If a tournament is not adjacent to a tournament with restriction (imposed by (CC) property) we distribute the probability uniformly
- The remaining tournaments (the middle layer) need a bit more work



## Theorem

There is a tournament rule that is monotone, (CC), (2-SNM- $\frac{1}{3}$ ), and (3-SNM- $\frac{1}{2}$ ).

- Previously, there was no known rule that was both (CC) and (3-SNM- $\alpha$ ) for $\alpha<\frac{31}{60}$.


## Rule extensions

- It is feasible to design fair and non-manipulable rules for a small number of teams on a computer
- We would like to extend these rules to more teams


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## Rule extensions

vertices of

- a manipulating group

- This extended rule is still (CC)
- If the size of a manipulating group is a lot smaller than $n$ then this rule is ( $k$-SNM- $(\alpha+\varepsilon)$ ) if the original rule is ( $k$-SNM- $(\alpha)$ )


## References

[1] Altman, A., Kleinberg, R.: Nonmanipulable randomized tournament selections. Proceedings of the National Conference on Artificial Intelligence 2, 686-690 (01 2010)
[2] Schvartzman, A., Weinberg, S.M., Zlatin, E., Zuo, A.: Approximately strategyproof tournament rules: on large manipulating sets and cover-consistence. In: 11th Innovations in Theoretical Computer Science Conference, LIPIcs. Leibniz Int. Proc. Inform., vol. 151, pp. Art. No. 3, 25, Schloss Dagstuhl. Leibniz-Zent. Inform., Wadern (2020)

## Thank you for your attention!

